

COL 100 - Lec 9

- Linear recursion
- Nested recursion
- Mutual recursion

$$f(n) = \begin{cases} 0 \\ f(n-1) + g(n-1) \end{cases}$$
$$g(n) = \begin{cases} 0 \\ g(n-1) + f(n) \end{cases}$$

i) $n=0$

otherwise

i) $n=0$

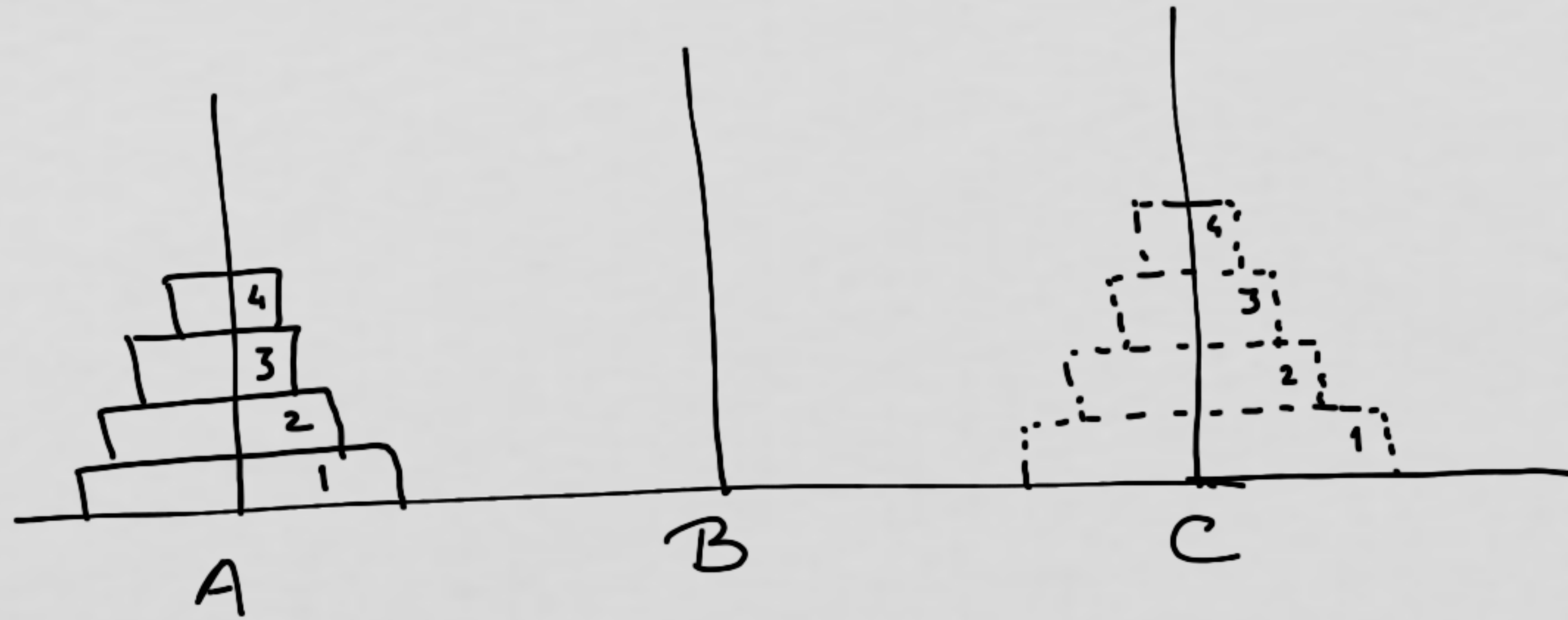
otherwise

$$g(g(2))$$
$$= g(g(1) + f(2))$$

Diagram illustrating the evaluation of $g(g(2))$ using mutual recursion:

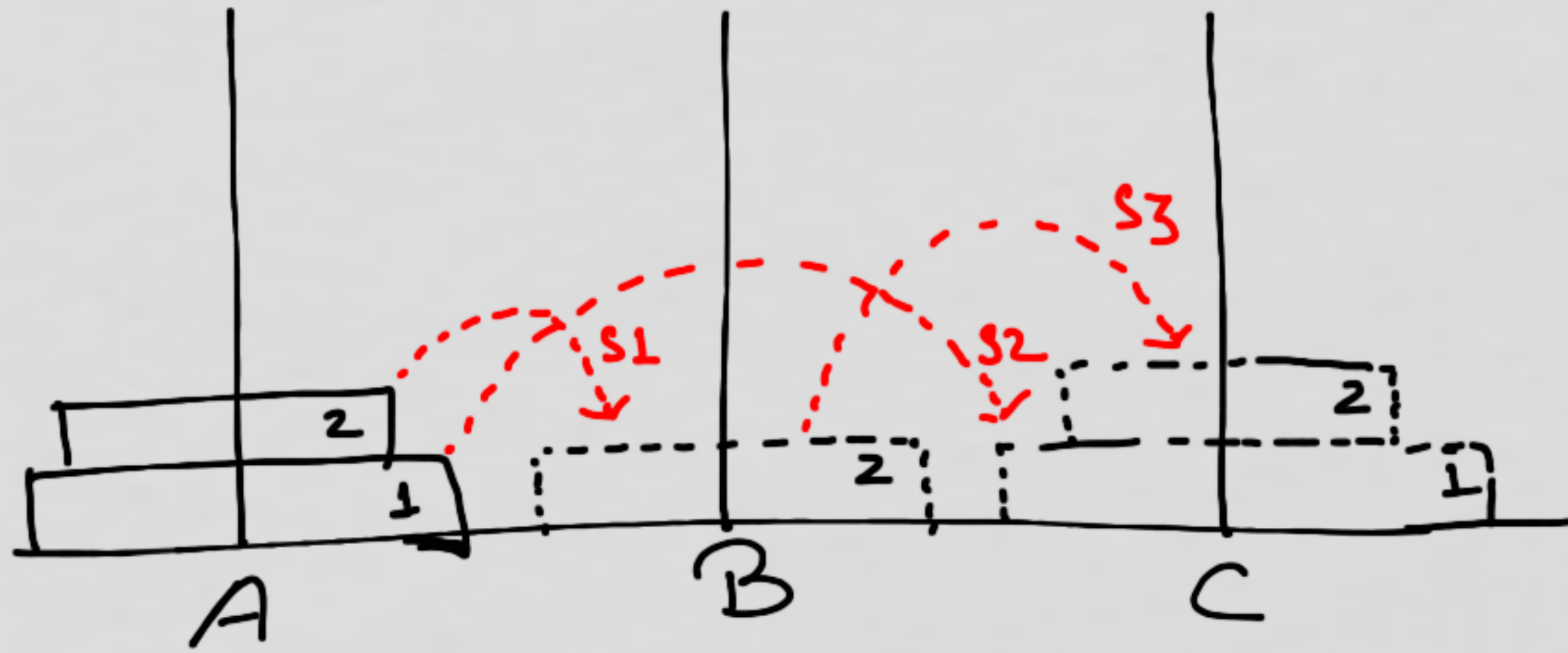
- $g(g(2))$ evaluates to $g(g(1) + f(2))$. The value 0 in $g(1)$ is highlighted in a red box.
- $g(g(1) + f(2))$ evaluates to $g(0 + f(1))$. The value 0 in $g(0)$ is highlighted in a red box.
- $g(0 + f(1))$ evaluates to $f(0) + g(0)$.
- $f(0) + g(0)$ evaluates to $f(1) + g(1)$.

Tower of Hanoi [Tower of Brahma]



Constraint : At no point a smaller disk
is beneath a larger disk!

Example with 2 disks



I have n disks, what observation can I make?

- Move top $(n-1)$ disks from A to B
- Move the n^{th} disk [bottom-most] from A to C
- Move $(n-1)$ disks from C to A

MoveDisks (n , src , aux , dst)

$=$ $\left\{ \begin{array}{l} \text{Move disk from src to dst} \\ \text{MoveDisks}(n-1, src, dst, aux) \text{ otherwise} \\ \text{Move disk from src to dst} \\ \text{MoveDisks}(n-1, aux, src, dst) \end{array} \right.$

i) $n=1$

Proof of Correctness

Base case: $n=1$, true

I.H. : $\exists n \geq 1 : \text{MoveDisks}(n, A, B, C)$ is true

I.S. : $\text{MoveDisks}(n+1, A, B, C) =$

- $\text{MoveDisks}(n, A, C, B)$ by I.H.
- Move n^{th} disk from A to C
- $\text{MoveDisks}(n, B, A, C)$

Recurrence Relation

$$T(n) = \begin{cases} 1 \\ 1 + 2T(n-1) \end{cases}$$

if $n = 1$

otherwise

$$T(n) = 2 * T(n-1) + 1$$

$$= 2^2 \cdot T(n-2) + 2 + 1$$

$$= 2^3 \cdot T(n-3) + 2^2 + 2 + 1$$

$$= 2^{n-1} T(1) + 2^{n-2} + \dots + 2^2 + 2 + 1$$

$$= 2^n - 1$$

Technical Completeness

fun pow(x, n) =

by default it is $\mathbb{N} \rightarrow \mathbb{N}$

• If I want to specify pow for ^{all} reals
then

```
fun pow(x: real, n: int) =  
  if n = 0 then 1.0  
  else pow(x, n-1) * x
```

Is it technically complete?

- Does your function guarantee, for instance,
 - $x \neq 0$ X
 - x is real ✓
 - n is int ✓
 - n is nonnegative X

```

fun pow(x: real, n) =
  if n < 0 then 1.0 / pow(x, ~n)
  else if n = 0 then 1.0
        else x * pow(x, n-1)
  
```

Is it technically complete? $0 \cdot 0^0 = 1.0$
 whereas $0 \cdot 0^n = 0 \ \forall n > 0$

• if $x = 0.0$
 and $n = -m < 0$

then $0.0^n = 1.0 / (0.0^m)$

$= \frac{1.0}{0.0}$

Division by zero

Pow int. version

$$x^n = \begin{cases} \text{undefined} \\ \text{undefined} \\ 1 \\ x \cdot x^{n-1} \end{cases}$$

if $n < 0$
 if $x = 0 \wedge n = 0$
 if $x \neq 0 \wedge n = 0$
 otherwise

Am1

Exception negExponent;

Exception ZeroPowerZero;

fun int pow (x, n) =

if n < 0

then raise negExponent

else if n = 0 then if x = 0 then

raise ZeroPowerZero

else 1

else x * intpow (x, n-1)