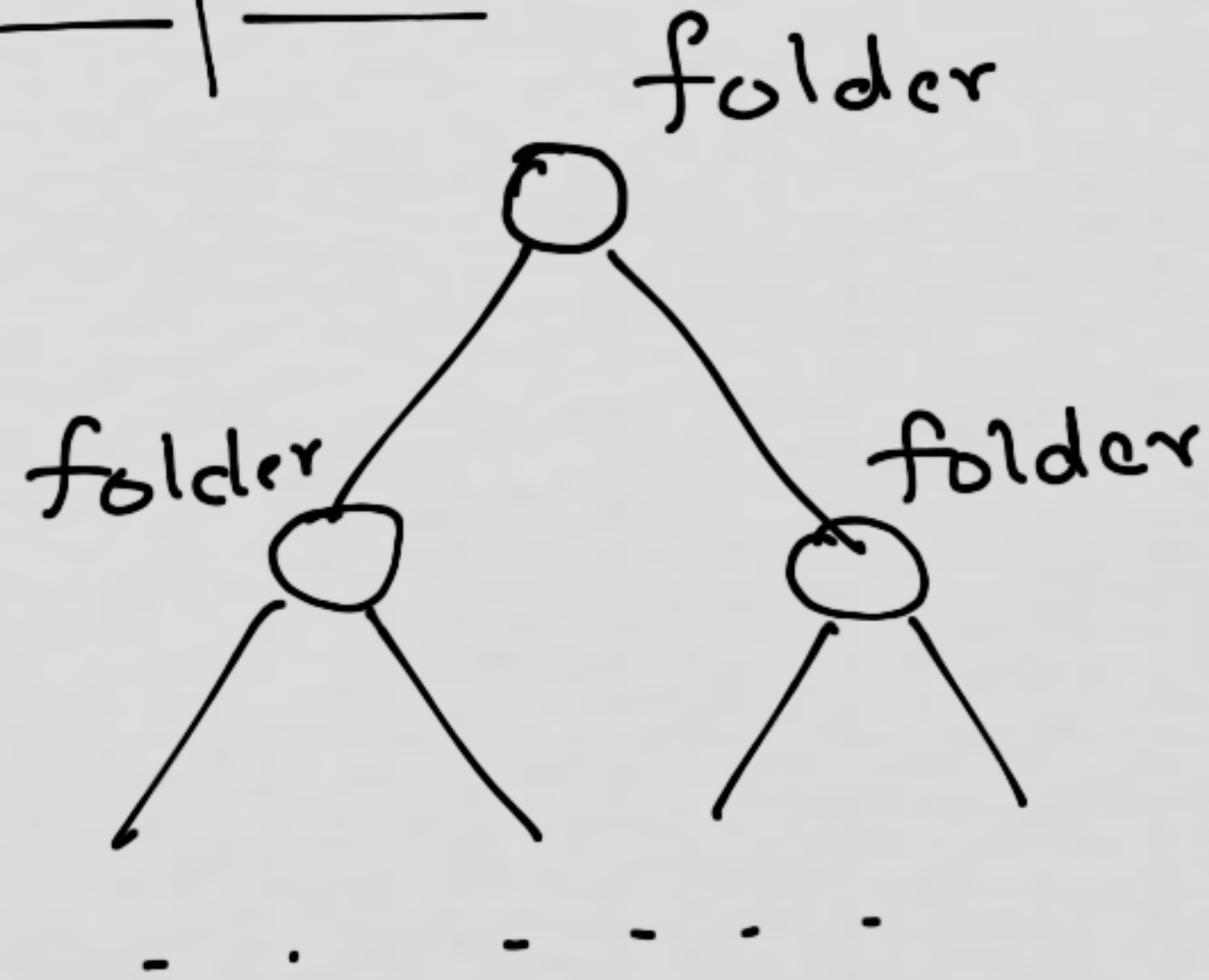


COL 100 - Lec 10

1. Uses of examples worked out
- Counting change → In vending machines
— more generally, it is a special case
of integer knapsack problem
[resource allocation]
 - Tower of Hanoi → disk backups,
neuropsychological tests etc.

More Examples



Searching for a folder
in the file system.

- get left child
- get Right child

Search (root, name Of Folder)

=

{

root

Not Found

Search (get left child (root),
name Of folder)

Search (get Right child (root),
name Of folder)

i | root's name
< = name Of folder

i | root's name
< = NULL

Otherwise

getLeftChild (node) = $\begin{cases} \text{NULL} \\ 1^{\text{st}} \text{ child} \end{cases}$

getRightChild (node) = $\begin{cases} \text{NULL} \\ 2^{\text{nd}} \text{ child} \end{cases}$

i) node is a
leaf
Otherwise

i) node is
a leaf

i) # of
Children = 2

Q: How does the above function unfold?
[Also called Depth first Search]

Q: Consider a binary tree, perform search
over a binary tree? Can you do better?

Ex 2:

Consider a 4 bit binary numbers

— find all 4-bit (in general N-bit) numbers without consecutive 1's.

i.e.

0000

0001

0010

0100

0101

1000

1001

1010

Base Case?

i) $n = 0$ then 0

[Base case not finished yet!]

Observations

- if last digit is 0
 - we can have both 0 & 1 in the current digit.

$$\text{CountNum}(n-1, 0) + \text{Count}(n-1, 1)$$

- if last digit is 1
 - we can only have 0 at current position

$$\text{CountNum}(n-1, 0)$$

Let us work out an example

CountNum (5, 0) initial condition

CountNum (3, 0) + CountNum (2, 1)

00
01
10

CountNum (2, 0) + CountNum (1, 1)

CountNum (1, 0) + CountNum (1, 1)

CountNum (1, 0) → 2

CountNum (1, 1) → 1

CountNum (2, 0)

CountNum (1, 0) → 2

New base cases

- When $n = 1$, last-digit = 0 then return 2
- When $n = 1$, last digit = 1 then return 1

Therefore

$$\text{CountNum}(n, k) = \begin{cases} 0 \\ 1 \\ 2 \\ \text{CountNum}(n-1, 0) + \text{CountNum}(n-1, 1) \\ \text{CountNum}(n-1, 0) \end{cases}$$

$$\text{if } n=0$$

$$\text{if } n=1 \wedge k=1$$

$$\text{if } n=1 \wedge k=0$$

$$\text{if } n>0 \wedge k=0$$

$$\text{if } n>0 \wedge k=1$$

Alternate solution

$n=1$	$n=2$	$n=3$	$n=4$	
0	00	000	0000	1000
1	01	001	0001	1001
	10	010	0010	1010
2	11	011	0011	1011
	3	100	0100	1100
		101	0101	1101
		110	0110	1110
		111	0111	1111
		5	5	3

X

Therefore
 $\forall n \geq 3, \text{Count}(N) = \text{Count}(N-1) + \text{Count}(N-2)$
 $\text{Count}(1) = 2, \text{Count}(2) = 3$

I I had asked you to generate
such binary strings, then?

GenNum(N , numStr, step, last-digit)

→ first generate all strings starting with '0'

- if $N = \text{step}$
 ↓
 then print numStr

- if last-digit is 1
 ↓
 then numStr ^ "0"

GenNum(N , numStr, $n+1$, 0)

- if last-digit is 0
 ↓
 then

GenNum (N, numStr[^]"0", n+1, 0)
GenNum (N, numStr[^]"1", n+1, 1)



String Concatenation Syntax

- May change from one language to other
- In C $\rightarrow$ string.concat (a, b)
- In SML $\rightarrow$ "Hello" ^ "_World"
- In python $\rightarrow$ "Hello" + "World"